Introduction to Braver Mohin Obstruction.

S Motivation. Q: Given y= 321 + 2. Find a rational solution for the equation Idea; Try  $\alpha = \frac{\alpha}{b} y = \frac{\alpha}{d}$  such that max [10], [b], [c], [d]] < B 'a constant. "herght" keep searching by chlarge B. If there is such rational solution. we can find one by this method. If there is no rational solution, we can't determine it by searching. Gi: How to determine solution set is empty ? More generally. for a ring R. we want to determine when the scheme.  $\chi = V(f_1, f_2, \dots, f_r) = \{(a_1, a_2, \dots, a_n) \in \mathbb{R}^n | f_1 = f_2 = \dots = f_n = 0\}$ has a solution in R. where fi is a polynomial with coefficient R. ie f; ER[x, ··· Xn]

Notation X(R) = {R solutions of X ]

= of R-valued points of Xy = Hom ( Speek, X). =  $Hom(R[x_1, ..., x_n]), R)$ ( $F_1 ... F_n$ )

Assume X is a finite type scheme over ring R. for the whole talk. I an algorithm to determine if  $X(R) = \phi$ King R Yes Ľ ies R Yes Fp Jes  $Q_{P}$ ? R Na ドル(け) No-2 What is Qp. (local field)?  $y^2 = 3a^3 t 2 \mod p$ , solution in  $\frac{y}{p_2}$ in 2/p24  $y^2 = 3a^3 + 2 \mod p^2$ .  $y^2 = 3a^3 + 2$  mod  $p^3$ in  $\mathbb{Z}_{p^{3}}\mathbb{Z}$ 

Take inverse limit. 
$$Z_p := \lim_{k \to \infty} Z_{p/2}^{k}$$
  
 $Q_p$  is the fractional field of integral domain  $Z_p$   
or:  $(Q_p = Z_p \partial_2 Q_p)$   
By construction, we know  
 $X(Q) \neq \emptyset \implies X(Q_p) \neq \emptyset$   
Generally. Given a number field  $k$ , and a prime  $V$   
we can construct a local field  $k_v$   
Similarly.  $X(k) = \phi \implies X(k) = \phi$   
Now we know  
 $X(k) \subset X(k_v)$  we have algorithm to determ  
 $X(k) \subset X(k_v)$  if  $X(k_v) = 0$ .  
So  $X(k) \subset TX(k_v)$   
Right hand side is a product of infinite term. It seems to be  
not computable anymore. We take a subset of it.  
Adelic point of  $X : X(A_k) \subset TX(k_v)$   
 $k_{k} := T'(k_v, O_v) = f(a_v) \in T_v k_v / \# fv(a_v \notin O_v f < \infty)$ 

If 
$$\chi$$
 is smooth projective geometrically integral variety  
then  $\chi(A_k) = T(\chi(K_k), \chi(O_k))$   
Now  $\chi(K) \longrightarrow \chi(A_k) \leftarrow$  we have algorithm  
to determine  $\chi(A_k) = \emptyset$   
what if  $\chi(A_k) \neq \emptyset$  how can we determine if  $\chi(K) = \emptyset$ ?  
Idea:  
Find a set T. such that  $\chi(K) \subset T \subset \chi(A_k)$   
such that  $\exists$  algorithm to determine if  $T = \emptyset$   
Even if  $\chi(A_k) \neq \emptyset$ . if we determine  $T = \emptyset$ , then we can  
also determine  $\chi(K) = \emptyset$   
T is called the obstruction  
That:  $\chi(K) \subset \chi(A_k)^{Br} \subset \chi(A_k)$   
 $\chi(A_k)^{Br}$  is called Braver - Mann obstruction.  
 $\int Braver groups$   
 $Br(k) = H_{ev}^2(Speck, Gm)$   
 $= H^2(G_k, (k^{Segr}))$  where  $G_k^2(Gall k^{Segr})$ 

= H (4k, (k.)) min un une = ? central simple algebras overky

Let A be a finite dimensional k-adjubra.  
A is called simple if A doesn't have non-trivial too side idea.  
A is called central if center of A is k.  
Central simple algebra of A over k if A is finite dimensional.  
k algebra that is central and simple.  
EX. Matrix algebra 
$$M_n(k)$$
 is central simple algebra.  
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EX. Halmitton's quaternion algebra H over IR is central simple  
algebra.  
[ Central simple algebra over K.]  
Define equivalent relation.  
 $A \sim B$  if  $A \otimes_k M_n(k) \cong B \otimes_k M_n(k)$  for some integer  
 $m n$ .  
Define Multiplication.  
 $[A] \times [B] = [A \otimes_k B]$   
 $A^{oPP}$  is the opposite algebra of A. i.e.  $\begin{cases} A^{oP} = A \text{ as a set (vector space)} \\ mdgitudion in A! X Y = ya \end{cases}$   
 $A \otimes A^{oPP} = M_n(k)$ .  
 $So [A] hos inverse.$   
 $Br(k) = l central simple algebra over k! 
 $So [A] hos inverse.$   
 $Br(k) = l central simple algebra over k! 
 $M = A \otimes_k M_n(k) = M_n(k)$  is a group.$$ 

Bauer groups of schemes.  
Dof: 
$$Br(X) = He_t^2(X, Gm)$$
  
Remarks it is also possible to generalize the third definition.  
Dof: An Azumaya algebra is a cohenent  $O_X$  - algebra A such that for any point  $\alpha \in X$ , the fitter  $A \otimes_{O_X} h(\omega)$  is a central simple algebra over the resolve field here.  
Two Azumaya algebra are similar to each other if there are locally free sheaf  $\mathcal{E}_1$  and  $\mathcal{E}_2$  such that  
 $A_1 \otimes_{O_X} End\mathcal{E}_1 \cong A_2 \otimes_{O_X} End\mathcal{E}_2$   
The similarity class of Azumaya algebra forms a group  
if X is regular and quasi - projective then the group is iso to  $Br(X)$   
 $\mathcal{S}$  Brauer - Manin obstruction.  
•  $Br(-) = H_{ec}^2(-, Gm)$  is a contravariant functor  
i.e.  $f: X \longrightarrow Y$  morphism of scheme.  
then it induced  $f^*: Br(Y) \longrightarrow Br(X)$   
 $X(K) \longrightarrow X(K_X)$   
Homic Speck X) Homic Speck, X).  
 $d \longrightarrow d_V$ 







X(Ax)<sup>Br</sup> consists of elements in X(Ax) or they onal to Br(X)  $\chi(k) \subset \chi(A_{k})^{Br} \subset \chi(A_{k})$ Q: Is it useful. In theory: Yes. There are many examples such that  $X(A_k) \neq \phi$ ,  $X(A_k)^{B_r} = \phi$ . We can use Brower Manin obstruction to argue that X(k)=0In fact: typically expect  $X(A_k)^A \neq X(A_k)$  "unless forced ortherwse" In practice: Not quite Yes. Can be computed in several examples but no genereral effective way to compute it. Assume X is a smooth projective variety over number fled K. now. Fact: If  $X(A_k)^{Br} = \phi$ , there exists a finite set  $B \subset Br(A_k)$ such that  $X(A_k)^{B} = \phi$ idea: Find subgp BC Br(X) such that  $X(A_k)^B$  is computable Def. We say B captures the Braver-Manin obstruction  $H X(A_{k})^{B_{i}} = \phi \implies X(A_{k})^{B} = \phi$ Theorem (DI) C is a smooth projective degree d genus / curve, then Br (C [d<sup>og</sup>]) (completely) captures the Braver-Manin obstruction obstruction. (2) If X is a smooth projective cubic obstruction then Br XI3] (completely) captures the Braver-Manin Obstruction.

For most example in literature, they show 
$$X(A_{1}^{k})^{B_{1}} \neq \phi$$
 by show  $X(A_{1}^{k})^{B_{2}} \neq \phi$  by show  $X(A_{1}^{k})^{B_{2}} \neq \phi$  by show  $X(A_{1}^{k})^{A_{1}} = 0$  for only one elements  $A \in B_{1}(X)$   
But generally, we need a bot elements in  $B_{1}(X)$ .  
Fact: Let  $A_{1}, A_{2} \in B_{1}(X)$  then  
 $X(A_{1})^{A_{1}} \cap X(A_{k})^{A_{2}} = \bigcap X(A_{k})^{T}$   $S_{k}^{k} \neq i$   
 $X(A_{1})^{A_{1}} \cap X(A_{k})^{A_{2}} = \bigcap X(A_{k})^{T}$   $S_{k}^{k} \neq i$   
we just need juste the generators  $A_{1}$  of  $B_{1}(X)$ .  
and calculate  $B_{1}(X)^{A_{1}}$ , then take intersection  
Thm. Let  $N \ge 0$ , char  $(K) \ne 2 \le number$  jield, we don't care.  
 $\exists$  smooth projective geometrically integral variety  $X$  over global jield  
 $S_{1}t. X(A_{k})^{B_{1}} = \phi$  but  $Y$  subge  $B \subset B_{1}(X)$  generated  
by  $< N$  elements ,  $X(A_{k})^{B} \ne \phi$ .

Let  $\pi: X \rightarrow \text{speek}$  be the structure morphism.  $B_{ro}(X) = \text{Im}(\pi^*; B_r(k) - B_r(k))$  'constant Braner classes' Fact:  $X(A_k) \stackrel{B_{ro}(X)}{=} X(A_k)$  $X(A_k)^{B_r}$  only depends on  $\frac{B_r(X)}{B_{ro}(X)}$ 

Hochschild-serve spectral sequence to Galois cover  $\overline{X} \rightarrow X$  and sheaf  $\operatorname{H}^{P}(G_{K}, \operatorname{Her}(\overline{X}, G_{m})) \rightarrow \operatorname{Her}(X, G_{m})$ 

exact sequence of low degree terms  

$$D \rightarrow Pic X \rightarrow (Pic\bar{X})^{G_{K}} \rightarrow Br K \rightarrow ker (Br(X) \rightarrow Br(\bar{X}))$$
  
 $\rightarrow H'(G_{K}, Pic\bar{X}) \rightarrow H^{3}(G_{K}, \bar{K}^{*})$   
 $\eta \to H'(G_{K}, Pic\bar{X}) \rightarrow H^{3}(G_{K}, \bar{K}^{*})$   
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Gabois modues, We can compute Br(X)/ Bro(X).

Reference: Rational points on varieties and the Brauer-Manin Ibstruction — Bianca Viruy The Brauer-Manin Obstruction — Shelly Manber